

Modeling Structured Event Streams in System Level Performance Analysis

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LCTES 2010, Stockholm, Sweden

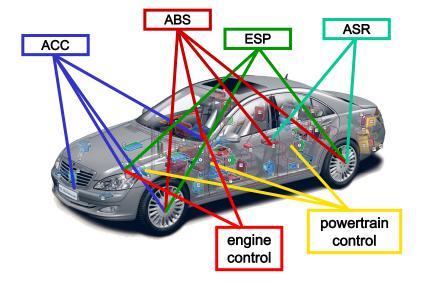


Outline

- Motivation
- Modular Performance Analysis with Real-Time Calculus
- Approach 1: FIFO Scheduling
- Approach 2: Event Count Curves
- Comparison / Case Study
- Conclusion

Motivation

Distributed Stream-based Embedded Systems



Design and analysis are complex due to:

- Concurrency
- Interference on shared resources
- Non-determinism in timing

Requirements for Performance Analysis:

- Coverage of system behaviors
- Accuracy of results
- Computational efficiency

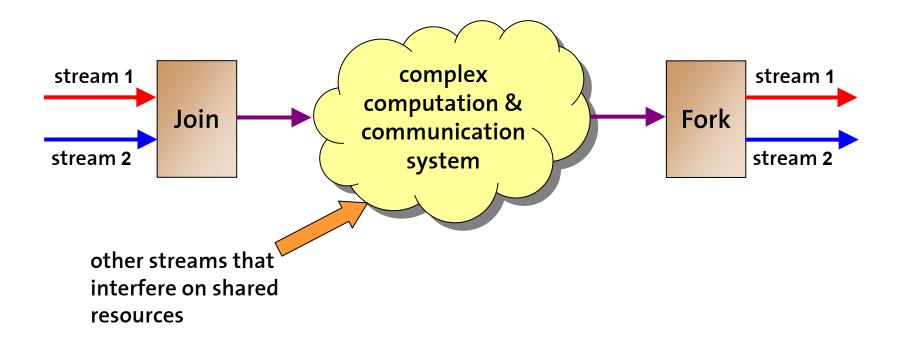
Analytic methods for Modular Performance Analysis: • MPA-RTC (ETHZ) • SymTA/S (TUBS)

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Motivation

Relevant design pattern: Join and Fork of event or data streams

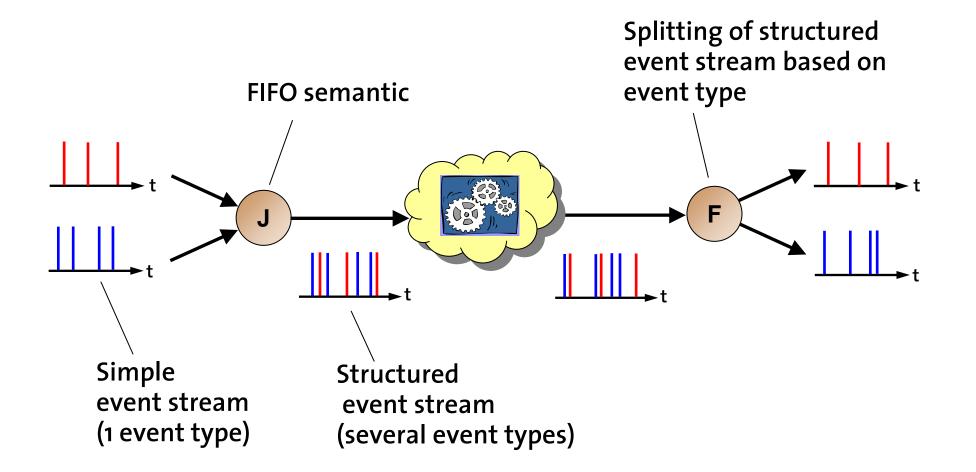


Goal: Extend analytic methods for performance analysis such that join/fork operations on streams can be captured accurately

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Definitions

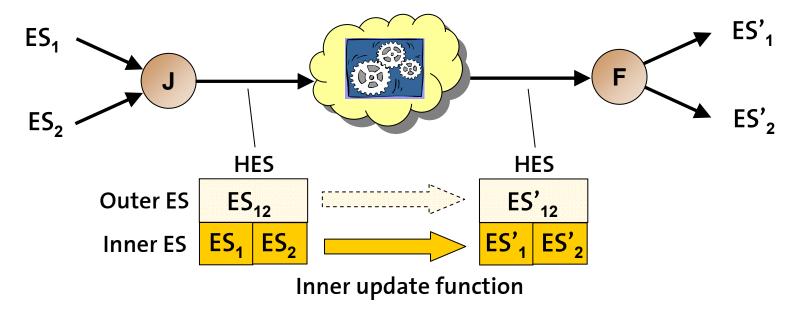


Event type = provenance

Related Work

Hierarchical Event Model (TU Braunschweig)

J. Rox & R. Ernst, *Modeling Event Stream Hierarchies with Hierarchical Event Models*, DATE 2008



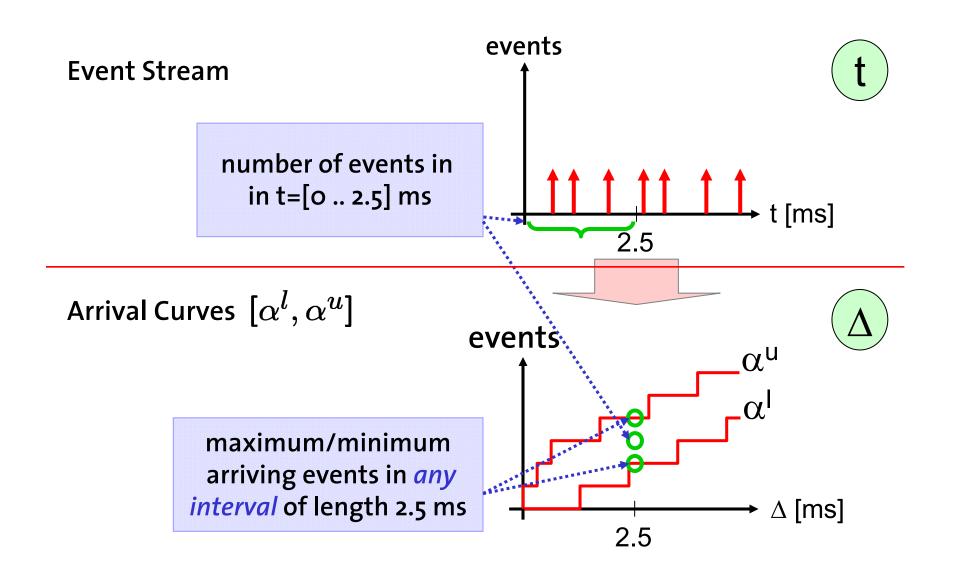
Limitations: - Method is not transparent to existing component models

- Deep processing of HES required by component models

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Arrival Curves



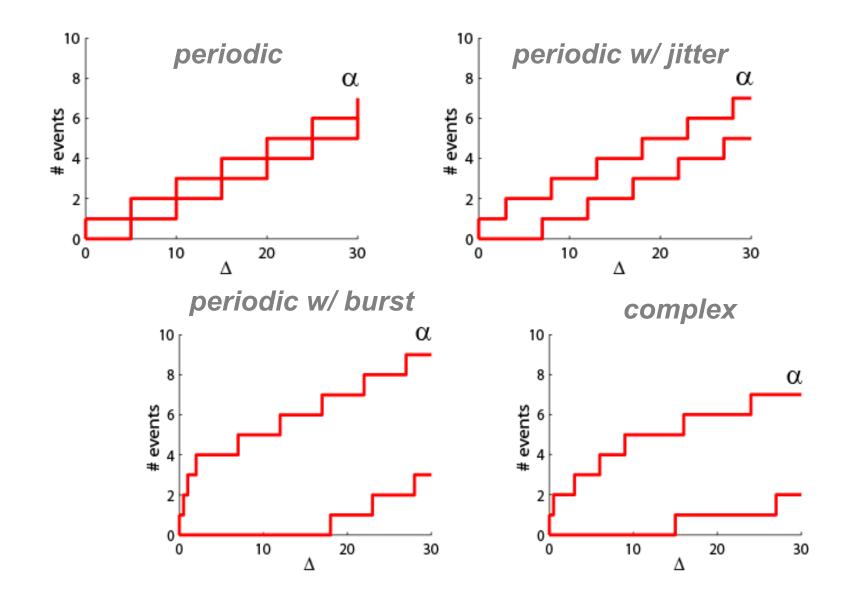
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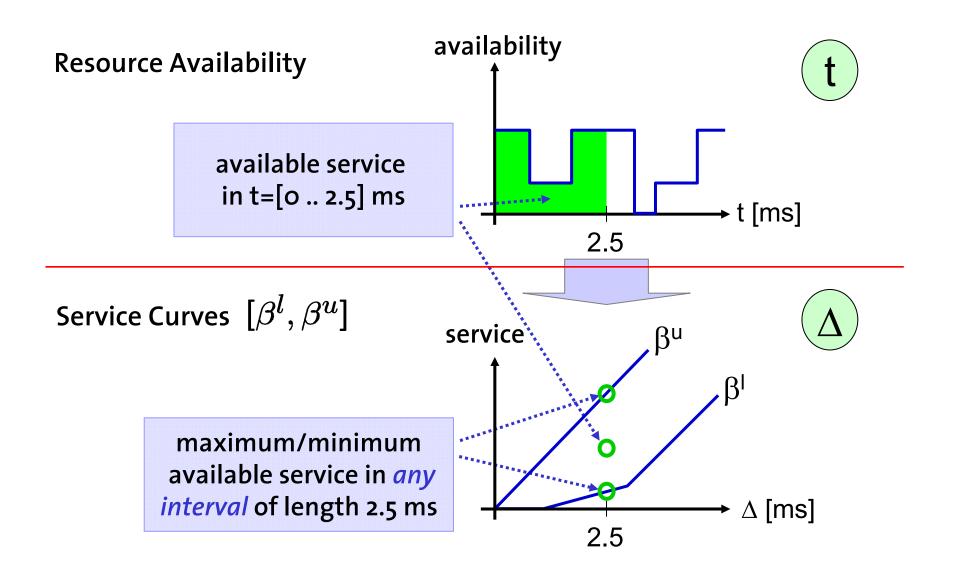
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Examples of Arrival Curves

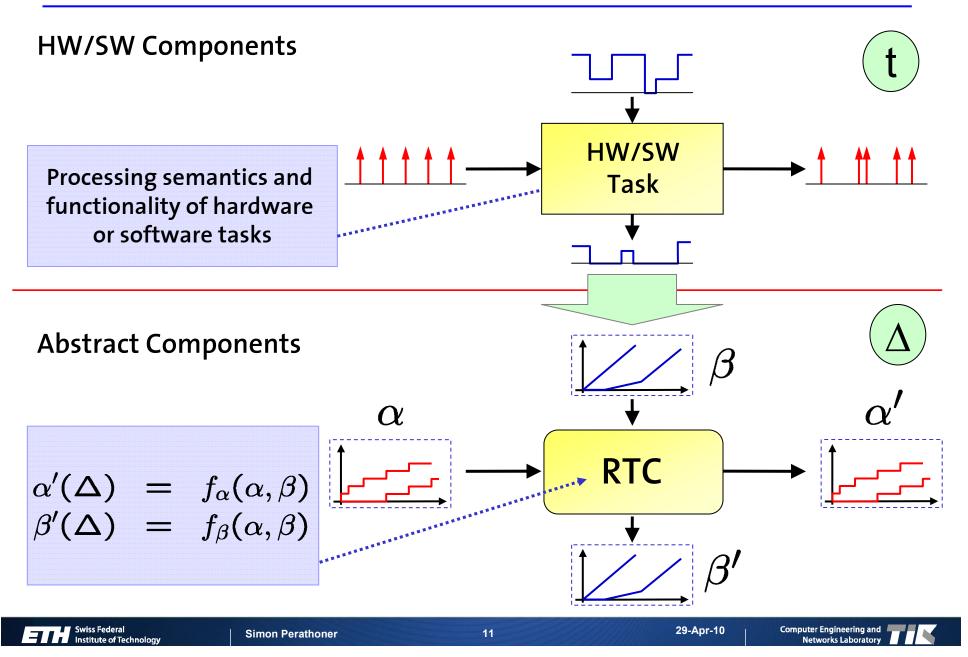


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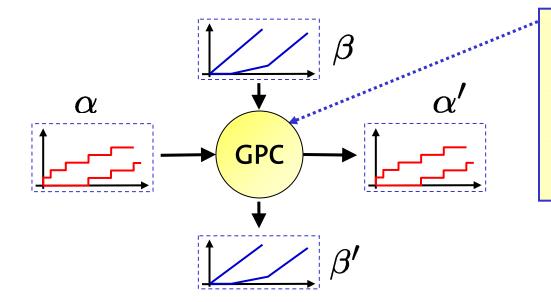
Service Curves



Processing Model (HW/SW)

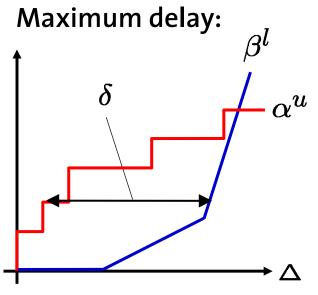


Example: Greedy Processing Component (GPC)



$$\begin{aligned} \alpha^{'u} &= GPC_{\alpha'^{u}}(\alpha^{u}, \beta^{u}, \beta^{l}) \\ \alpha'^{l} &= GPC_{\alpha'^{l}}(\alpha^{l}, \beta^{u}, \beta^{l}) \\ \beta'^{u} &= GPC_{\beta'^{u}}(\beta^{u}, \alpha^{l}) \\ \beta'^{l} &= GPC_{\beta'^{l}}(\beta^{l}, \alpha^{u}) \end{aligned}$$

- FIFO input buffer
- Greedy event processing



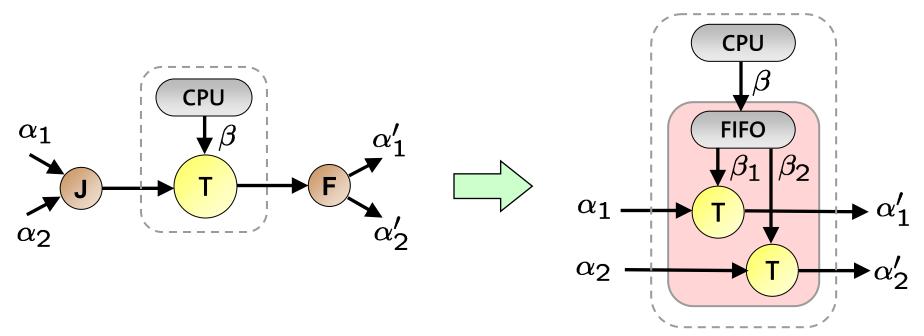
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FIFO Scheduling

- Keep sub-streams separated in the model
- Adapt existing component models



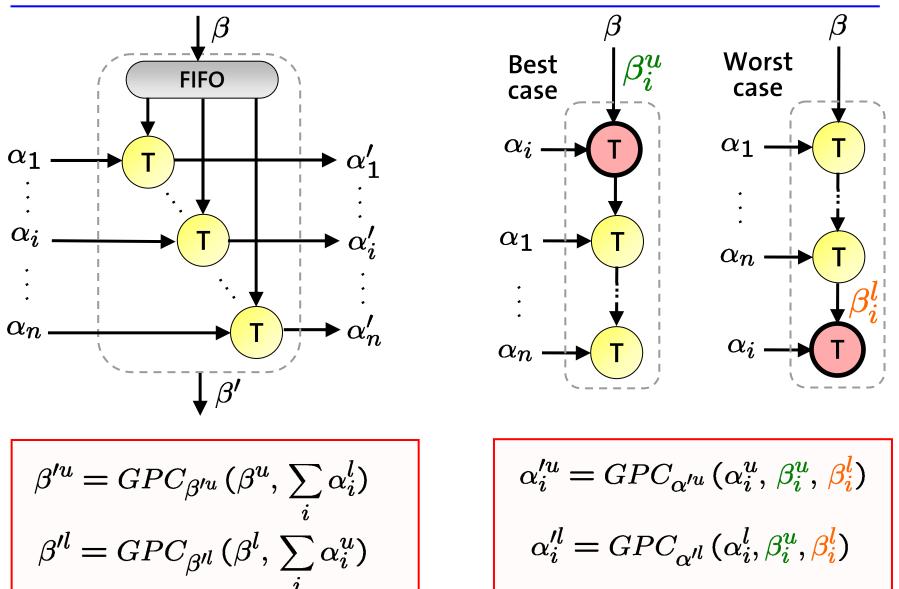


Modeling approach justified by FIFO semantics of Join operator

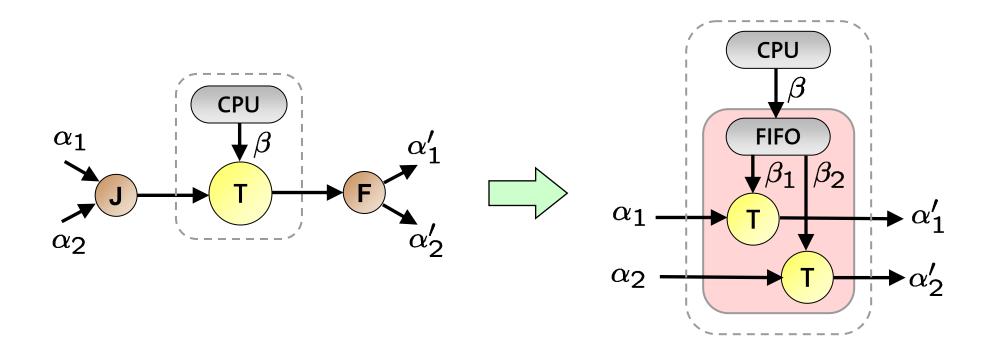
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FIFO Component in MPA



FIFO Scheduling



Abstract processing component handles structure of input stream explicitly

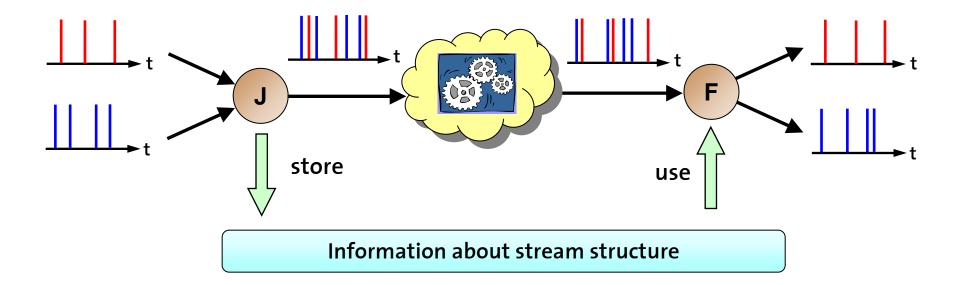
 \Rightarrow Method is not transparent to existing component models!

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Event Count Curves: Basic Idea

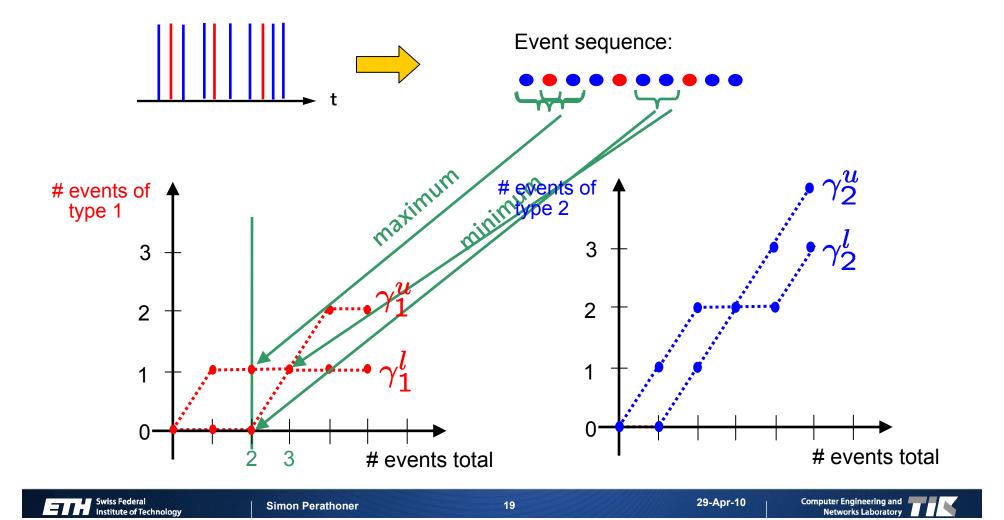
The processing components in MPA preserve the order of the events in a stream



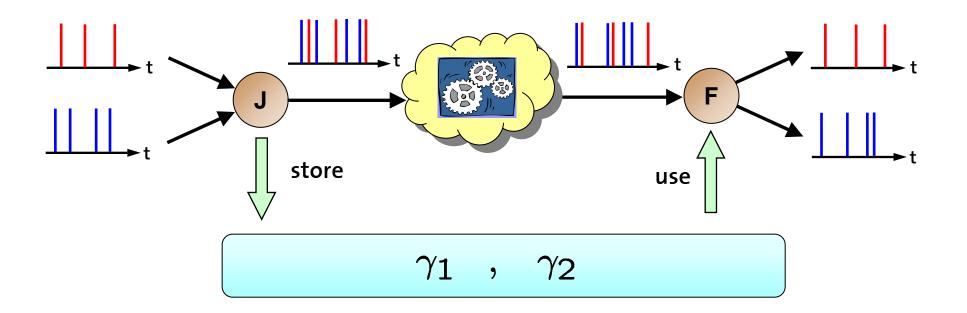


Event Count Curves: Definition

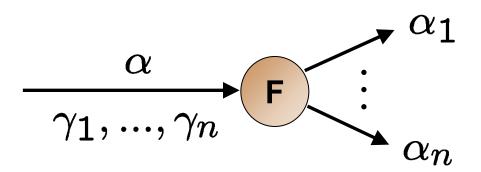
For a fixed number of events in the structured stream the lower (upper) ECC $\gamma_i^l(\gamma_i^u)$ bounds the min (max) number of events of type i.



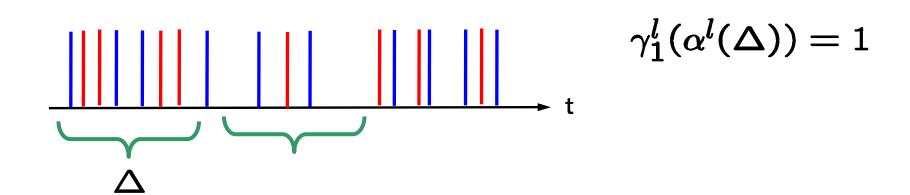
Event Count Curves



Fork (Simple Event Streams)

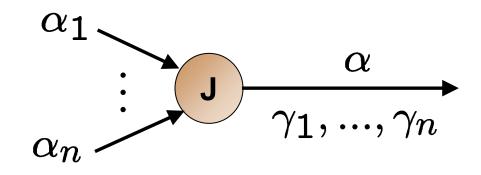


$$\alpha_i^l(\Delta) = \gamma_i^l(\alpha^l(\Delta))$$
$$\alpha_i^u(\Delta) = \gamma_i^u(\alpha^u(\Delta))$$



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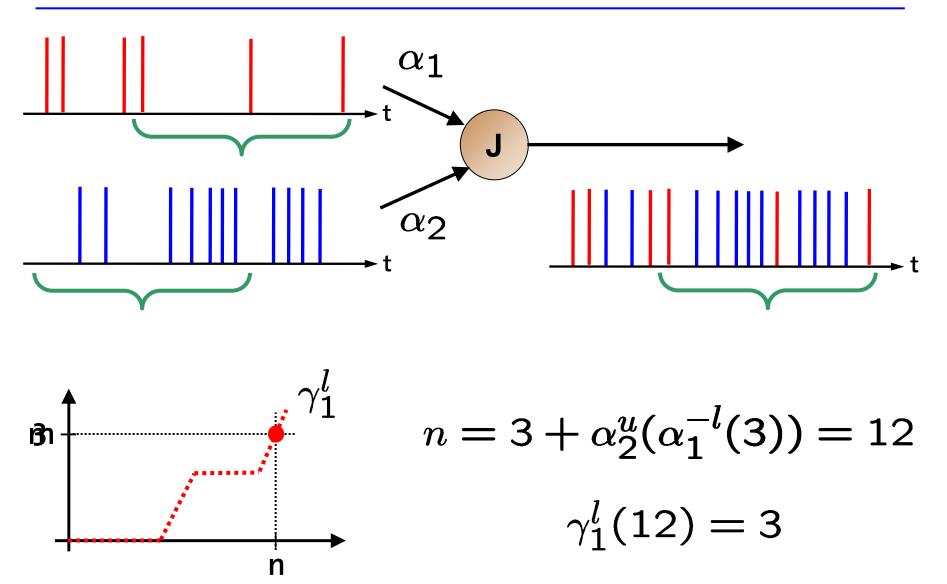
Join (Simple Event Streams)



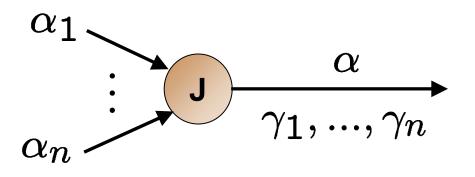
$$lpha^l(\Delta) = \sum_i lpha_i^l$$
 ; $lpha^u(\Delta) = \sum_i lpha_i^u$



Join (Simple Event Streams)



Join (Simple Event Streams)



$$lpha^l(\Delta) = \sum_i lpha_i^l$$
 ; $lpha^u(\Delta) = \sum_i lpha_i^u$

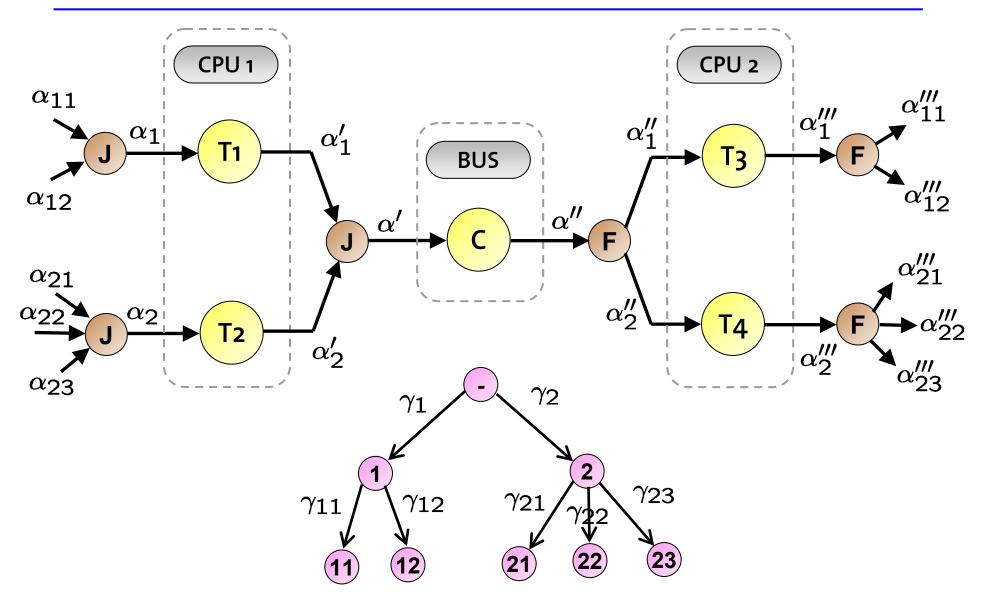
$$\gamma_i^l(n) = \epsilon_i^{-u}(n) \qquad \epsilon_i^u(n_i) = n_i + \sum_{j \neq i} \alpha_j^u(\alpha_i^{-l}(n_i))$$

$$\gamma_i^u(n) = \epsilon_i^{-l}(n) \qquad \epsilon_i^l(n_i) = n_i + \sum_{j \neq i} \alpha_j^l(\alpha_i^{-u}(n_i))$$

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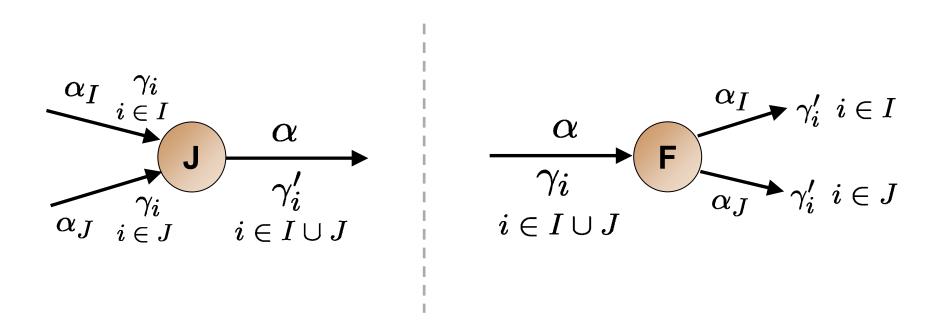
Hierarchical application of ECCs



Arbitrary Join and Fork of Structured Streams

Limitation of hierarchical ECC organization:

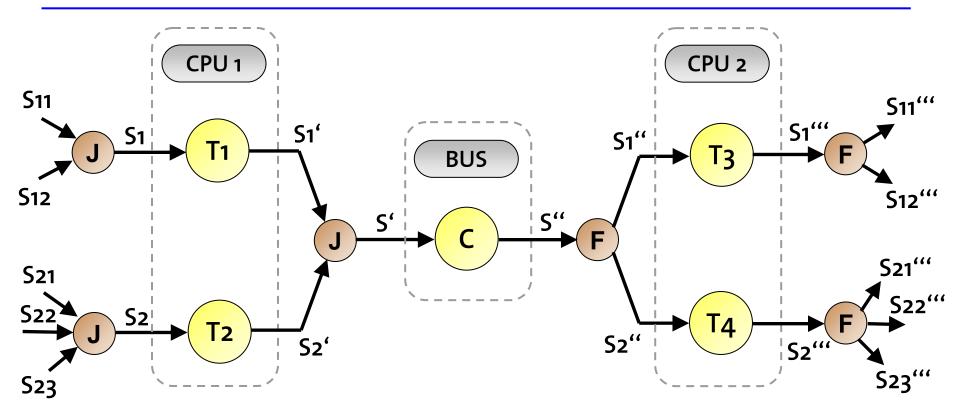
- Structured streams can be decomposed only in the way they have been composed!
- ⇒ More general join/fork operators are desirable



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Example

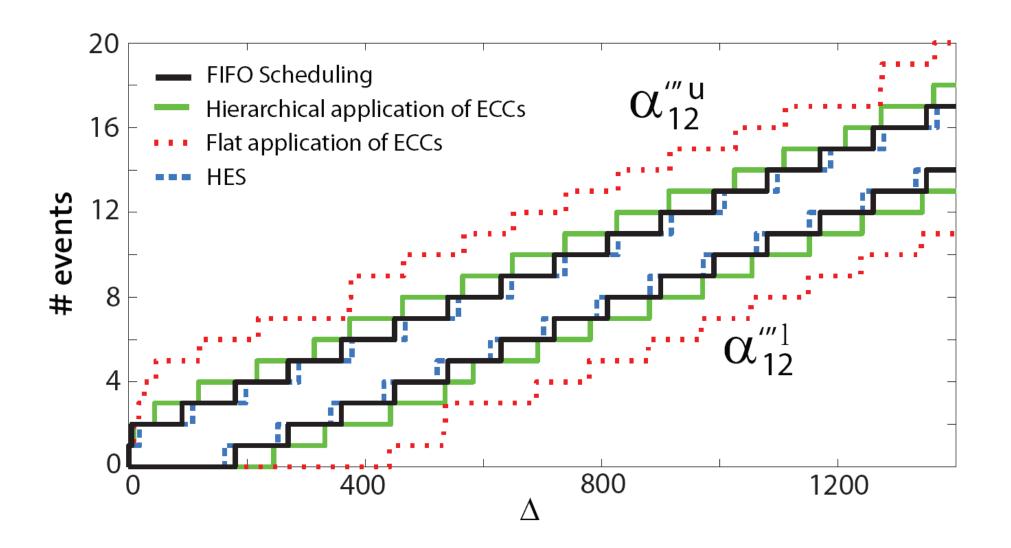


Stream	Period	Jitter		Task	BCET	WCET
S11	100	30	-	T1	2	3
S12	90	15		T2	3	4
S21	30	о		C	9	12
S22	80	20		T3	2	3
S23	75	5		T4	4	5

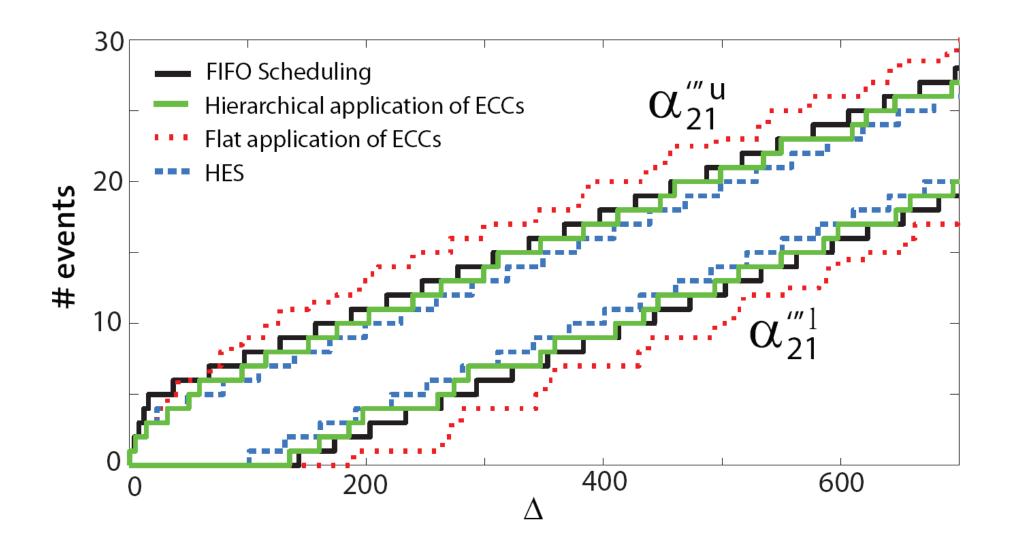
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Analysis Results

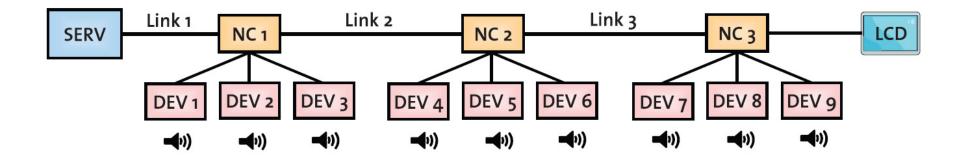


Analysis Results



Case Study: Specification

Distributed information/entertainment application



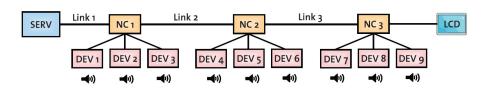
2 classes of network traffic:

• On-demand audio streaming (high prior.)

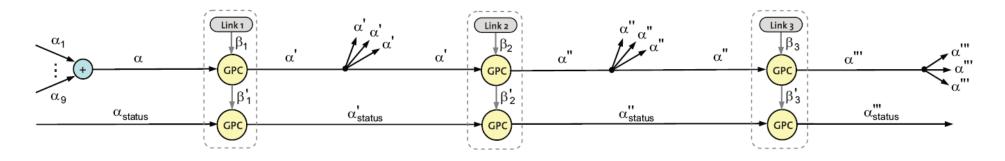
• Real-time status information (low prior.)

	Size	Period	Jitter	Deadline
Audio Frames	1'518 Bytes	30 ms	5 ms	100 ms
Status Frames	106'500 Bytes	5 s	0 s	1.5 s

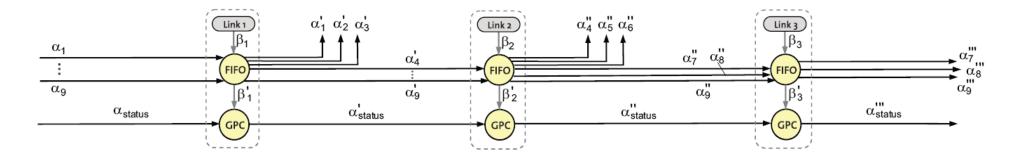
Case Study: Models (1)



Classic MPA

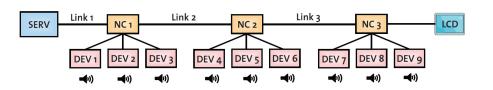


FIFO components

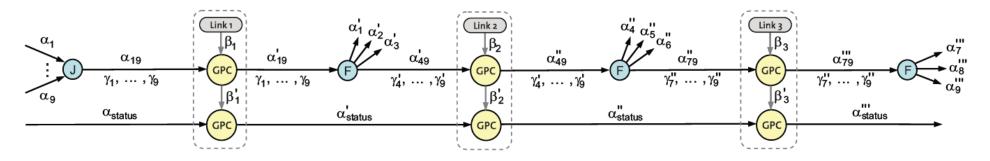


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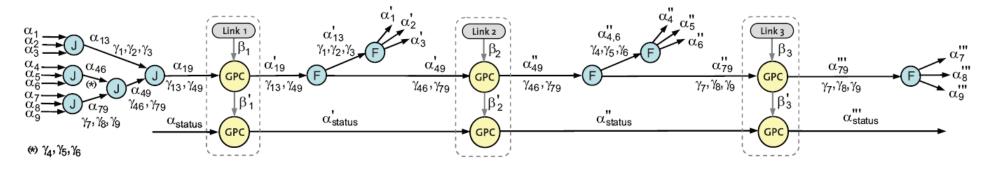
Case Study: Models (2)



ECC (flat)

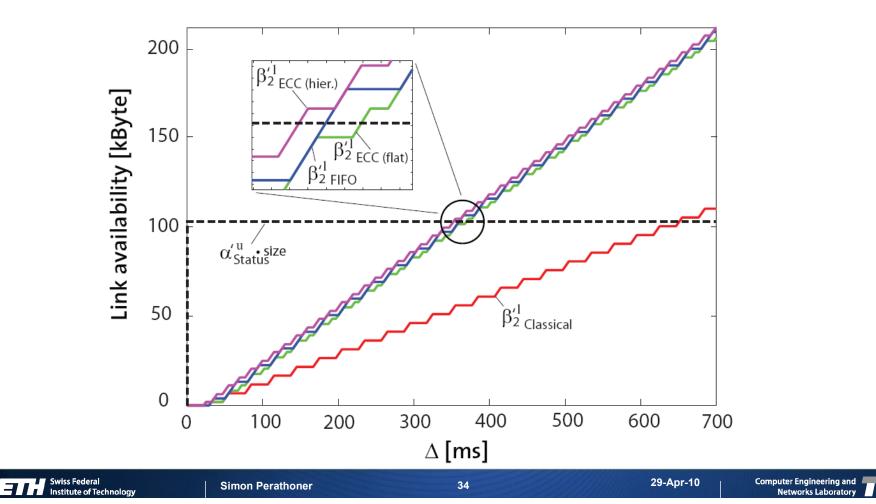


ECC (hierarchic)



Case Study: Results

	Classic	FIFO	ECC (flat)	ECC (hierar.)
Max. delay	1.954 s	1.255 s	1.316 s	1.248 s



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Conclusion

- Two new methods for modeling and analysis of joined event streams in modular performance analysis
- The concept of Event Count Curves is orthogonal to existing event stream models (PJD, arrival curves) and transparent to all analysis components ⇒ Compositionality of methods is not affected
- ECC operators for structured streams allow *arbitrary* decomposition of streams
- Comparisons and Case Study highlight the utility of the proposed methods



Thank you!

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